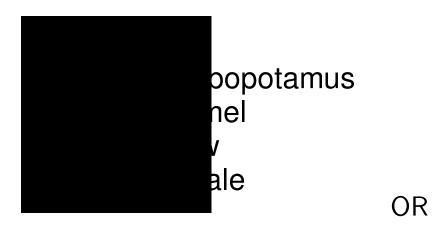
#### Announcements

- Homework 1 is due Tuesday November 9 22:00 discussion regarding questions in MS Teams
- Work on the journal club

(read the paper, plan the meeting no later than Nov. 23)

**Evolution and Phylogenetic Trees** 

Broňa Brejová October 28, 2021





# Phylogenetic tree reconstruction (fylogenetický strom)

# Input:

# m aligned sequences, each of length n

human	С	А	G	Т	Т	А
elf	А	А	Т	А	G	А
Gollum	С	С	G	А	G	А
hobbit	С	С	G	Т	Т	С
orc	А	А	Т	Т	Т	А

# **Output:**

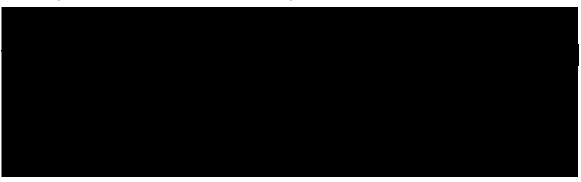
tree representing their evolutionary history



Newick format: (((gollum,hobbit),human),(elf,orc))

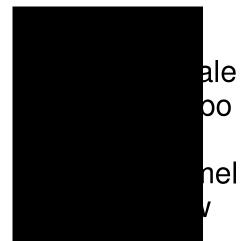
#### Rooted and unrooted trees

Unrooted tree (nezakorenený strom)



Two out of seven possible rooted versions of the tree





Most methods reconstruct unrooted trees

# Rooting a tree using an outgroup

Add outgroup (dog) to the unrooted tree





# Parsimony principle and maximum parsimony (úspornosť)

**Input:** (aligned) sequences of several extant species.

**Task:** Find a phylogenetic tree that explains the data by using the **minimum number of evolutionary events**.

Here: Evolutionary event = single base mutation

**Subtask:** For a given phylogenetic tree, find **ancestral sequences** that require the minimum number of events (score of the tree)

gollum	AGA	
hobbit	TAA	
human	TCC	
elf	ACA	
orc	TCA	





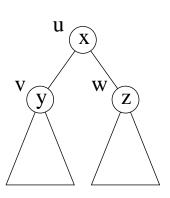
# Computing cost of a given phylogenetic tree

Use dynamic programming (separately for each alignment column).

For each internal vertex u and symbol x:

 $N_{u,x}$ : how many events are required in the subtree of u, assuming that the symbol in u is x?

 $N_{u,x} = \min_{y} \{ N_{v,y} + [x \neq y] \} + \min_{z} \{ N_{w,z} + [x \neq z] \}$ 





**Time:** O(m)Repeat for each alignment column: O(mn)

#### What we have: compute the cost of a particular tree

gollum	AGA
hobbit	TAA
human	TCC
elf	ACA
orc	TCA





5 changes

#### What we want: Find the tree with the smallest cost

gollum	AGA
hobbit	TAA
human	TCC
elf	ACA
orc	TCA



#### Finding the most parsimonious tree

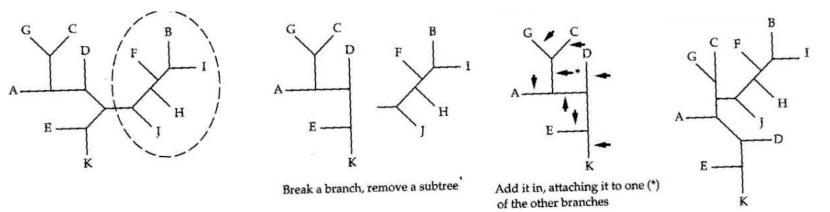
# **NP-hard problem**

**Trivial algorithm:** try all possible trees.

For m species  $1 \cdot 3 \cdot 5 \cdots (2m - 5) = (2m - 5)!!$ E.g. for 10 species cca 2 mil., for 20 species  $2 \cdot 10^{20}$ 

# Heuristic search:

- Start with some "sensible" tree
- Explore similar trees by using e.g. "subtree pruning and regraft":



# Neighbour joining (metóda spájania susedov)

- We throw away "details" of which mutations happened
- Summarize by a **distance matrix**  $D_{ij}$

#### **Example:**

human	C	٨	C	т	т	۸		hu	е	h	ho	0
							human	0	4	3	2	2
elf	A	A	Т	A	G	A	elf	4	0	3	6	2
gollum	С	С	G	А	G	А						
hobbit	С	С	G	Т	Т	С	gollum					
orc	Δ	Δ	т	т	т	Δ	hobbit	2	6	3	0	4
or c	7 (	7 (	•	·	I	7 (	orc	2	2	5	4	0

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## Idea of neighbour joining

Assume that the distances  $D_{i,j}$  correspond to the real distances in the tree (they are **additive**)

		gollum	hobbit	human	elf	orc
	gollum	0	5	9	15	16
ł	hobbit	5	0	8	14	15
ł	human	9	8	0	16	17
e	elf	15	14	16	0	3
ç o	orc	16	15	17	3	0

 $D_{\text{hobbit,human}} = 2 + 1 + 5 = 8$ 

#### Idea of neighbour joining

- Assume that the distances  $D_{i,j}$  correspond to the real distances in the tree (they are **additive**)
- Find two leaves *i* and *j*, for which we can **say with certainty**, that they have the same parent in the tree
- Join i and j and replace them with a parent node k with new distances to each other node  $\ell$ :

$$D_{k,\ell} = \frac{D_{i,\ell} + D_{j,\ell} - D_{i,j}}{2}$$

#### How to find out which two leaves should be joined?



Choose leaves i, j minimizing:

$$L_{i,j} = (m-2)D_{i,j} - \underbrace{\sum_{k \neq i} D_{i,k}}_{r_i} - \underbrace{\sum_{k \neq j} D_{j,k}}_{r_j}$$

m: the number of leaves

Connect leaves i, j, which minimize the following quantity:

$L_{i,j} = (m-2)D_{i,j} -$	– $\sum D_{i,k}$ –	$-\sum D_{j,k}$
	$k{ eq}i$	$k{ eq}j$
	$r_i$	$r_j$

	D		L	new $D$
	g ho hu e o	$r_i$	g ho hu e o	g ho hu e+o
g	0 5 9 15 16	45 g	72 -68 -58 -48	g 0 5 9 14
ho	5 0 8 14 15	42 hc	-7268 -48 -48	ho 50813
hu	9 8 0 16 17	50 հւ	ı -68 -6850 -50	hu 98015
е	15 14 16 0 3	48 e	-58 -48 -50 . <b>-90</b>	e+o 14 13 15 0
0	16 15 17 3 0	51 o	-48 -48 -50 <b>-90</b> .	

**Running time of neighbor joining:**  $O(m^3)$  (m: number of leaves) In 2009 a  $O(m^2)$  version was developed (Elias and Lagergren)

#### Neighbour joining: summary

- If the distance matrix is additive and corresponds to the real evolutionary distances then neighbour joining finds the correct tree
- Longer sequences  $\Rightarrow$  better distance estimates  $\Rightarrow$  correct trees
- How to compute "real" evolutionary distances? Counting differences is not enough

hna a n	C	۸	C	т	т	٨		hu	е	g	ho	0
human							human	0	4	3	2	2
elf	A	A	Т	A	G	A	elf	4	0	3	6	2
gollum	С	С	G	А	G	А	gollum					
hobbit	С	С	G	Т	Т	С	-					
orc	Δ	Δ	т	т	т	Δ	hobbit					
ore	Л	Л		ı	1		orc	2	2	5	4	0

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#### **Problems with estimating distances**

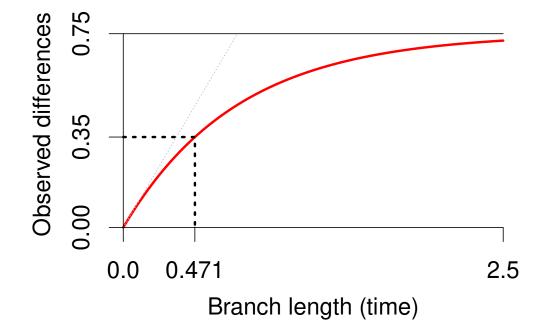
- One base may mutate multiple times during evolution (possibly even back to original base)
- When counting differences we see at most one change at each position ⇒ we underestimate the real distance
- We want a correction to estimate the real number of mutations that have occurred

#### Jukes-Cantor substitution model

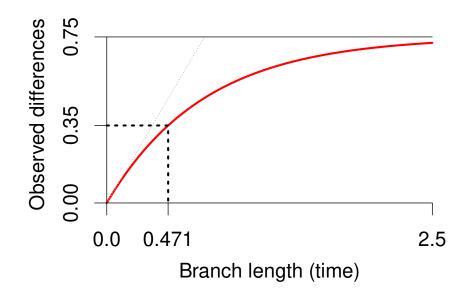
Probability that base A changes to C in time *t*:  $Pr(X_{t_0+t} = C | X_{t_0} = A) = \frac{1}{4}(1 - e^{-\frac{4}{3}\alpha t})$ 

 $\alpha$ : mutation rate (the number of substitutions per unit of time)

Expected number of observed changes per base in time *t*:  $D(t) = \Pr(X_{t_0+t} \neq X_{t_0}) = \frac{3}{4}(1 - e^{-\frac{4}{3}\alpha t})$ 



#### Back to distances in neighbor joining



• Using this model, we can correct observed distances

$$D = \frac{3}{4}(1 - e^{-\frac{4}{3}\alpha t}) \qquad \Rightarrow \qquad \alpha t = -\frac{3}{4}\ln(1 - \frac{4}{3}D)$$

• Next week: more complex models of evolution

# Maximum likelihood trees (najvierohodnejšie stromy)

A phylogenetic tree with branch lengths can be viewed as a **simple generative model** 

$$X_{root}$$

$$X_{gbm}$$

$$X_{gbm}$$

$$X_{gbm} = A$$

$$X_{gb}$$

$$X_{gb}$$

$$X_{eo}$$

$$X_{gb} = A$$

$$X_{eo} = C$$

$$X_{g} = X_{b}$$

$$X_{m}$$

$$X_{e}$$

$$X_{o}$$

$$X_{g} = A$$

$$X_{m} = G$$

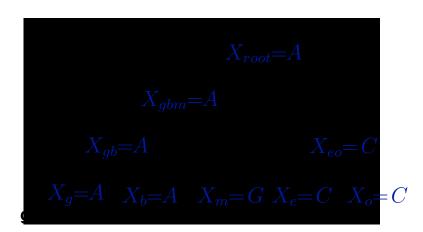
$$X_{e} = C$$

$$X_{g} = A$$

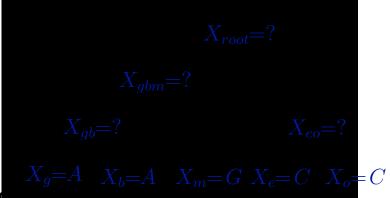
$$X_{m} = G$$

$$X_{e} = C$$

Probability that it generates particular bases in nodes:  $Pr(X_g = A, X_b = A, X_m = G, X_e = C, X_o = C, X_{gb} = A, X_{gbm} = A, X_{eo} = C, X_{root} = A)$   $= Pr(X_{root} = A) \cdot Pr(A \mid A, t_1) \cdot Pr(C \mid A, t_2) \cdot Pr(A \mid A, t_3) \cdot Pr(G \mid A, t_4) \cdot Pr(A \mid A, t_5) \cdot Pr(A \mid A, t_6) \cdot Pr(C \mid C, t_7) \cdot Pr(C \mid C, t_8)$   $Pr(C \mid A, t_2) \text{ is a abbreviation of } Pr(X_{eo} = C \mid X_{root} = A), \text{ J.-C. model}$  We can compute (product):



We want to compute **tree likelihood**:



# Likelihood of a tree (vierohodnosť stromu):

 $Pr(X_g = A, X_b = A, X_m = G, X_e = C, X_o = C)$ Add up probabilities of all letter combinations in ancestors  $X_{gb}$ ,  $X_{gbm}$ ,  $X_{eo}$ ,  $X_{root}$ 

Compute using Felsenstein algorithm

(simple dynamic programming similar to the parsimony)

For a given alignment, tree and branch lengths we can compute likelihood in O(nm) time

# How to find the tree with the highest likelihood?

- Again NP-hard problem ; complicated because we also need **branch lengths**
- Typical heuristic algorithm:
  - Start with a "reasonable" tree
  - Compute its likelihood
    - \* Start with "reasonable" branch lengths
    - \* Compute likelihood using these branch lengths
    - \* Iteratively improve branch lengths to improve the likelihood (e.g. gradient descent)
  - Explore "similar" trees to improve likelihood (as with parsimony).

# Consistency of algorithms for phylogeny

- "Well-behaved" algorithms: if the length of the sequences n increases, the answer should get closer to the correct answer.
- The algorithm for phylogeny is consistent, if the probability of obtaining the correct tree converges to 1 with n → ∞.

# Algorithm comparison

	Complexity	Consistency	Data utilization
Parsimony	NP-hard	NO	complete sequences
Neighbor Joining	$O(m^3)$	YES	distances only
Likelihood	NP-hard	YES	complete sequences

# Sources of data for phylogenetic trees

Some special sequences are often used (e.g. ribosomal RNA genes, mitochondrial genome)

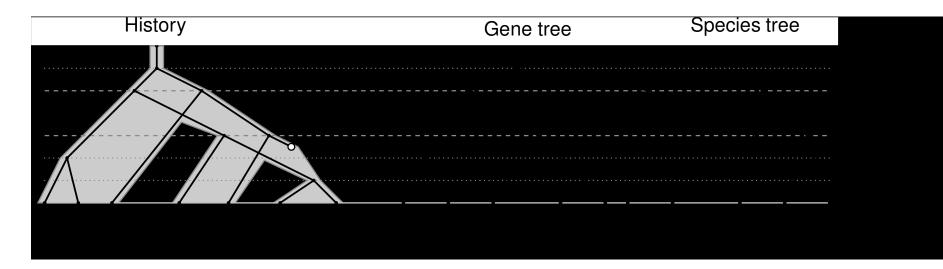
# What about using DNA sequences of other genes?

- Choose a suitable gene
- Find its homologs in other species
- Use these to construct the tree (DNA sequences or proteins)

Problem: genes can be duplicated and lost in evolution

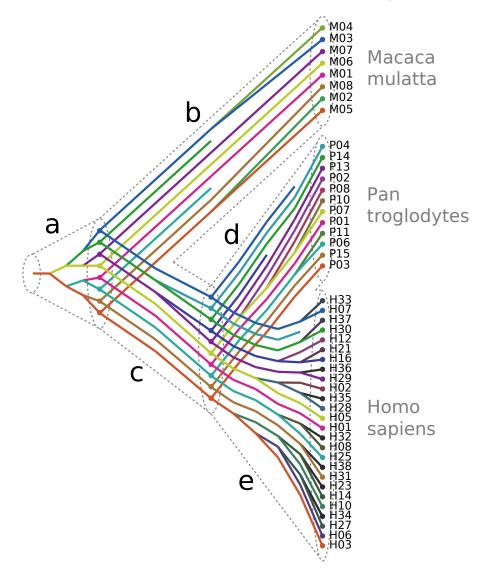
#### History of a duplicated gene

**Example:** species a, b, c, genes  $a_1, a_2, a_3, b_1, b_2, c_1, c_2$ 



- Homologs: similar sequences evolved from a common ancestor
- Orthologs: closest common ancestor is a speciation (e.g. pairs of genes a<sub>1</sub> - b<sub>1</sub>, a<sub>2</sub> - b<sub>1</sub>)
- Paralogs: closest common ancestor is a duplication (e.g. pairs of genes a<sub>1</sub> - a<sub>2</sub>, a<sub>1</sub> - b<sub>2</sub>)

# A more complex example of gene duplication:



# Summary

Substitution models allow us to:

- estimate real evolutionary distance (the number of substitutions) from the observed difference count between two sequences
- compute the probability that we observe a particular nucleotide change over time t

Three methods for phylogeny inference:

- Parsimony
- Neighbour joining
- Maximum likelihood

Gene trees and species trees, complications in phylogeny reconstruction